## Search-based Local Blackbox Deobfuscation: Understand, Improve and Mitigate

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## Speaker



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## Obfuscation



## Obfuscation



## Deobfuscation



## Deobfuscation

## Protecting Software through Obfuscation: Can It Keep Pace with Progress in Code Analysis?

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## A Generic Approach to Automatic Deobfuscation of Executable Code

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## Symbolic deobfuscation:

from virtualized code back to the original*

Jonathan Salwan ${ }^{1}$, Sébastien Bardin ${ }^{2}$, and Marie-Laure Potet ${ }^{3}$

Backward-Bounded DSE: Targeting Infeasibility Questions on Obfuscated Codes*
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## Deobfuscation

Protecting Software through Obfuscation: Can It Keep Pace with Progress in Code Analysis?

Backward-Bounded DSE: Targeting Infeasibility Questions on Obfuscated Codes*

Babak Yadegari
Brian Johannesı

Whitebox deobfuscation is highly efficient
o the original*

## Whitebox Deobfuscation

## But efficient countermeasures

Information Hiding in Software with Mixed<br>Boolean-Arithmetic Transforms

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How to Kill Symbolic Deobfuscation for Free (or: Unleashing the Potential of Path-Oriented Protections)

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Probabilistic Obfuscation through Covert Channels

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## New threat: Blackbox Deobfuscation



Syntia: Synthesizing the Semantics of Obfuscated Code
Tim Blazytko, Moritz Contag, Cornelius Aschermann, and Thorsten Holz, Ruhr-Universität Bochum
https://www.usenix.org/conference/usenixsecurity 17/technical-sessions/presentation/blazytko

This paper is included in the Proceedings of the 26th USENIX Security Symposium August 16-18, 2017 • Vancouver, BC, Canada

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## Bypasses whitebox methods limitations

## Open questions



## Contributions



Propose missing formalization

Refine Syntia experiments: new strengths and weaknesses

Show and explain why MCTS is not appropriate


## Mitigate



S-metaheuristics > MCTS
Implement our approach:
Xyntia

Evaluation of Xyntia

Propose 2 protections

Evaluate them against Xyntia and Syntia

## The talk in a nutshell

## I. Blackbox deobfuscation: what's that?

II. Deepen understanding
III. Improve state-of-the art IV. Mitigate


## Blackbox deobfuscation : what's that ?

## Blackbox deobfuscation

## 1) Sample

$$
\begin{aligned}
& (t=1, T=2) \\
& (t=2, T=5) \\
& (t=0, T=6)
\end{aligned}
$$


-1
-3
-6
$\cdot$

## 2) Learn

$$
\begin{aligned}
& (\mathrm{t}=1, \mathrm{~T}=2) \rightarrow-1 \\
& (\mathrm{t}=2, \mathrm{~T}=5) \rightarrow-3 \\
& (\mathrm{t}=0, \mathrm{~T}=6) \rightarrow-6
\end{aligned}
$$


$t-T$

## Learning engine

$$
\begin{aligned}
& U+(T-1) \quad t+T \quad t-U \\
& U \times U \quad(t-T) \times(T-1)
\end{aligned}
$$



## Expression Grammar

$$
\begin{aligned}
U:= & U+U|U-U| U * U \ldots \\
& |t| T \mid 1
\end{aligned}
$$

## Why blackbox?

## Given a language $L$ and an expression " $e$ " in $L$

## Syntactic complexity

Size of the the expression "e"
Size of the smallest expression in $L$ equivalent to "e"

## Example

$t-T$ is syntactically simpler than $(t \vee-2 T) \times 2-(t \oplus-2 T)+T$ but they share the same semantic complexity (being equivalent)

## Why blackbox ?

## Given a language $L$ and an expression " $e$ " in $L$

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Obfuscation increase syntactic complexity
$\rightarrow$ No impact on blackbox methods

## Understand

## Zoom on SoA: Syntia



- Dig into Syntia and deepen its evaluation:
- RQ1: stability of Syntia
- RQ2: efficiency of Syntia
- RQ3: Impact of operators set


## Syntia: new results



## Correctness



## Syntia: new results



## Quality



## Correctness

## Robustness



## Experimental design

## B1 (Syntia)

- 500 expressions
- Use up to 3 inputs
- redundancy
- Unbalanced w.r.t. type


## B2 (ours)

- 1110 expressions
- Use 2-6 inputs
- No redundancy
- Balanced w.r.t. type

|  | Type |  |  |  |  | \# Inputs |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Bool. | Arith. | MBA |  | 2 | 3 | 4 | 5 | 6 |  |
| \#Expr. | 370 | 370 | 370 |  | 150 | 600 | 180 | 90 | 90 |  |

## Table 1: Distribution of samples in benchmark B2

## Evaluation of Syntia

## B1 (Syntia)

- With a 60 s/expr. timeout : 75\% of success rate
- With a $1 \mathrm{~h} /$ expr. timeout : $88.2 \%$ of success rate
- With a 12 h/expr. timeout : 97.6 \% of success rate


## B2 (Ours)

Table 2: Syntia depending on the timeout per expression (B2)

|  | 1 s | 10s | 60s | 600 s |
| :---: | :---: | :---: | :---: | :---: |
| Succ. Rate | 16.5\% | 25.6\% | 34.5\% | 42.3\% |
| Equiv. Range | 16.3\% | 25.1-25.3\% | 33.7-34.0\% | 41.4-41.6\% |
| Mean Qual | 0.35 | 0.49 | 0.59 | 0.67 |

## Why ?



$$
(U+U \xrightarrow[U-U]{\Delta+U-U}
$$

入 ( Observed Samples:

$$
\begin{array}{ll|l|}
(\mathrm{t}=1, \mathrm{~T}=2) & \rightarrow & -1 \\
(\mathrm{t}=10, \mathrm{~T}=0) & \rightarrow & \mathbf{1 0} \\
(\mathrm{t}=10, \mathrm{~T}=5) & \mathbf{5} &
\end{array}
$$

## Why ?



$$
\begin{array}{cc}
U+U \quad U-U & T
\end{array} \stackrel{t}{t}
$$

$\Delta$

$$
\begin{aligned}
& \text { ( Observed Samples: } \\
& \begin{array}{ll|l}
(\mathrm{t}=1, \mathrm{~T}=2) & \rightarrow & \mathbf{- 1} \\
(\mathrm{t}=10, \mathrm{~T}=0) \rightarrow & \mathbf{1 0} \\
(\mathrm{t}=10, \mathrm{~T}=5) \rightarrow \\
5
\end{array},
\end{aligned}
$$

## Why ?

$$
\begin{array}{cc}
5+T-T & U+U \quad U-U
\end{array}
$$

1 Observed Samples:

$$
\begin{array}{ll|}
(\mathrm{t}=1, \mathrm{~T}=2) & -\mathbf{1} \\
(\mathrm{t}=10, \mathrm{~T}=0) \rightarrow & \mathbf{1 0} \\
(\mathrm{t}=10, \mathrm{~T}=5) \rightarrow & 5
\end{array}
$$

Synthesized Samples:

$$
\begin{aligned}
& (\mathrm{t}=1, \mathrm{~T}=2) \rightarrow \mathbf{0} \\
& \begin{array}{l}
(\mathrm{t}=10, \mathrm{~T}=0) \rightarrow-\mathbf{- 1 1} \\
(\mathrm{t}=10, \mathrm{~T}=5) \rightarrow-6
\end{array} \\
& (\mathrm{t}=10, \mathrm{~T}=5) \rightarrow-6
\end{aligned}
$$

## Why ?

## (C) $t-T$

$$
\begin{gathered}
U+U+U-U \\
(U+U)-U \frac{(U-U)-U}{2} \\
(T-t)-1
\end{gathered}
$$

入 Observed Samples:

| $(\mathrm{t}=1, \mathrm{~T}=2)$ | $-\mathbf{1}$ |
| :--- | :--- |
| $(\mathrm{t}=10, \mathrm{~T}=0) \rightarrow$ | $\mathbf{1 0}$ |
| $(\mathrm{t}=10, \mathrm{~T}=5) \rightarrow$ | 5 |

Synthesized Samples:

$$
\begin{array}{ll|c}
(\mathrm{t}=1, \mathrm{~T}=2) & \rightarrow \\
(\mathrm{t}=10, \mathrm{~T}=0) & \rightarrow \\
(\mathrm{-11} \\
(\mathrm{t}=10, \mathrm{~T}=5) \rightarrow-\mathbf{- 6}
\end{array}
$$

## Why ?



$$
\begin{array}{ccc}
U+U \rightarrow U-U & T & t \\
(U+U)-U & (U-U)-U & \vec{U}-t
\end{array}
$$

$\Delta$

$$
\left(\begin{array}{l}
\text { Observed Samples: } \\
(\mathrm{t}=1, \mathrm{~T}=2) \rightarrow \\
(\mathrm{t}=10, \mathrm{~T}=0) \rightarrow \mathbf{- 1} \\
(\mathrm{t}=10, \mathrm{~T}=5) \rightarrow \\
\mathbf{5}
\end{array}\right),
$$

## Why ?


$\Delta$


)

## Why ?

$$
t-T
$$

$$
\begin{aligned}
& U \\
& U+U \quad U-U \quad{ }_{T} \quad{ }_{t}^{*} \\
& (U+\vec{U})-U \quad(U-\vec{U})-U \quad \vec{U}-T \quad \vec{U}-t \\
& \text { そ } \zeta \\
& t-T \quad 1-T
\end{aligned}
$$

$\Delta$

$$
\left(\begin{array}{c}
\text { Observed Samples: } \\
(\mathrm{t}=1, \mathrm{~T}=2) \rightarrow \\
(\mathrm{t}=10, \mathrm{~T}=0) \rightarrow \mathbf{1 0} \\
(\mathrm{t}=10, \mathrm{~T}=5) \rightarrow \mathbf{5}
\end{array},\right.
$$

## Why ?

$$
\text { (6) } t-T
$$

$$
\begin{aligned}
& \underset{U+U}{U-U} \\
& (U+U)-U \quad(U-\vec{U})-U \quad \underset{~-~}{U-T} \\
& 1-T \\
& \text { © } \\
& \text { Synthesized Samples: } \\
& \begin{array}{ll}
(t=1, T=2) & \rightarrow \\
(t=10, T=0) & \rightarrow \mathbf{- 1} \\
(t=10, T=5) & \rightarrow \\
& \mathbf{- 4}
\end{array}
\end{aligned}
$$

## Why ?



## Claim

- Search space is too unstable for partial node evaluation
- Estimation of non terminal expressions is misleading

Evidence $\mathbf{n}^{\circ} 1$ : 2 simulations can lead to very distinct distances

Evidence n ${ }^{\circ} 2$ : Syntia does not benefit from partial evaluation

Evidence n ${ }^{\circ}$ 3: Syntia behaves in practice almost as BFS

## Evidence $\mathbf{n}^{\circ} 1$ and 2



Evidence $\mathbf{n}^{\circ} 2$

Number of expressions from exploitation steps?

## B1 (Syntia)

20 / 376

Non terminal expressions

## B2 (Ours)

34 / 341
Simulation leads to completely different results

## Evidence nº3

## - Config. of Syntia makes MCTS almost BFS

##  <br>  <br> Syntia is not guided $\uparrow$ <br> Over B2 Syntia and enum. MCTS reach similar results

## Improve 回

## Blackbox deobf., an optimization pb

Syntia sees blackbox deobfuscation as a single player game


We propose to see it as an optimization problem

Goal : find $\underset{\text { an expr. }}{s^{*}}$ s.t. $\underbrace{f}_{\Delta}\left(s^{*}\right) \leq f(s), \forall s \in S$

## S-metaheuristics

- Solve optimization problems



## S-metaheuristics

- Solve optimization problems



## S-metaheuristics

- Solve optimization problems



## S-metaheuristics

- Solve optimization problems



## S-metaheuristics

- Solve optimization problems



## S-metaheuristics

## - Solve optimization problems



## New prototype: Xyntia

## ) Xyntia

S-metaheuristics

Can choose between:
$\rightarrow$ Hill Climbing
$\rightarrow$ Simulated annealing
$\rightarrow$ Metropolis Hasting
$\rightarrow$ Iterated Local Search


MCTS

## Xyntia vs Syntia

## B1 (Syntia)

- $\mathbf{1 0 0}$ \% success rate in 1 s/expr.


B2 (Ours)


## Xyntia vs Syntia

## B1 (Syntia)

- 100 \% success rate in 1 s/expr.

B2 (Ours)


-     - Stable
- Good quality results


## Other experiments

- Xyntia against QSynth
- Xyntia against "compiler like simplifications"
- Xyntia against program synthesizer CVC4
- Xyntia against superoptimizer STOKE
- Use-cases:
- State-of-the-art protections
- VM-based obfuscation



## What's next?



## DOWHWORT1/

## DHITSACMIM

## Mitigate <br> 

## Context : Virtualization



## Proved to be sensitive to blackbox deobfuscation




## Why VM-based obf. is vulnerable ?



- Handlers are too semantically simple:
$\rightarrow$ e.g. $+,-, \times, \wedge, v$
- Obfuscation increase syntactic complexity $\rightarrow$ Blackbox deobf. is not impacted

We need to move ...
From syntactic to semantic complexity

## Semantically complex expressions

## - Goal:

- Increase the semantic complexity of each handlers
- Keep a Turing complete set of handlers
- Example:

$$
\begin{array}{ccc} 
& h_{0}= & (x+y)+-\left(\left(a-x^{2}\right)-(x y)\right) \\
+\quad & h_{1}=\left(a-x^{2}\right)-x y+(-(y-(a \wedge x)) \times(y \otimes x)) \\
+\quad h_{2} & = & (y-(a \wedge x)) \times(y \otimes x) \\
\hline & h= & x+y
\end{array}
$$

## Merged handlers

## - Goal:

- Increase semantic + sampling complexity


## - Example:

$$
\begin{gathered}
h_{1}(x, y)=x+y \quad \text { and } \quad h_{2}(x, y)=x \wedge y \\
\rightarrow \quad h(x, y, c)=\text { if }(c=c s t) \text { then } h_{1}(x, y) \text { else } h_{2}(x, y)
\end{gathered}
$$

## - Need to hide conditionals:

```
int32_t h(int32_t a, int32_t b, int32_t c) {
    // if (c == cst) then h1(a,b,c) else h2(a,b,c);
    int32_t res = c - cst ;
    int32_t s = res >> 31;
    res = (-((res ^ s) -s) >> 31) & 1;
    return h1(a, b, c)*(1 - res) + res*h2(a, b, c);
}
```


## Semantically complex handlers: results



## More results:

- Syntia with 12 h /exprs. $\rightarrow 1 / 15$ on BP1


## Merged handlers: results



Figure 10: Merged handlers: Xyntia (timeout=60s)

## More results:

- Syntia finds nothing for $\geq 2$ nested ITE


## Conclusion



MCTS is not appropriate for blackbox deobfuscation
$\rightarrow$ Search space too unstable
$\rightarrow$ Estimation of non terminal expressions pertinence is misleading

S-metaheuristics yields a significant improvement
$\rightarrow$ More robust
$\rightarrow$ Much Faster

Moving for syntactic to semantic complexity
$\rightarrow 2$ efficient methods to protect against blackbox deobfuscation

## Thank you for your attention



